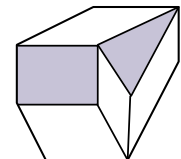

Some New Techniques for Explicit Mixed Volume Computation

Reinhard Steffens, Thorsten Theobald

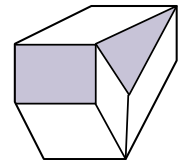


Erlangen, 19. September 2008



Mixed Volumes

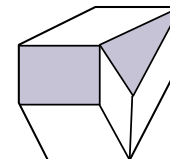
Let P_1, \dots, P_n be n polytopes in \mathbb{R}^n and $\lambda_1, \dots, \lambda_n$ non-negative real parameters. Then $\text{vol}_n(\lambda_1 P_1 + \dots + \lambda_n P_n)$ is a homogeneous polynomial of degree n in $\lambda_1, \dots, \lambda_n$ with non-negative coefficients. The coefficient of $\lambda_1 \cdots \lambda_n$ is called the **mixed volume** of P_1, \dots, P_n and is denoted by $\text{MV}_n(P_1, \dots, P_n)$.



Mixed Volumes

Let P_1, \dots, P_n be n polytopes in \mathbb{R}^n and $\lambda_1, \dots, \lambda_n$ non-negative real parameters. Then $\text{vol}_n(\lambda_1 P_1 + \dots + \lambda_n P_n)$ is a homogeneous polynomial of degree n in $\lambda_1, \dots, \lambda_n$ with non-negative coefficients. The coefficient of $\lambda_1 \cdots \lambda_n$ is called the **mixed volume** of P_1, \dots, P_n and is denoted by $\text{MV}_n(P_1, \dots, P_n)$.

$$\text{MV}_n(P_1, \dots, P_n) = \sum_{\substack{C \text{ mixed cell of a} \\ \text{mixed subdivision} \\ \text{of } P_1 + \dots + P_n}} \text{vol}_n(C)$$

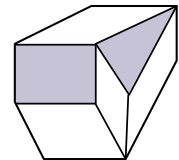


Mixed Volumes

Let P_1, \dots, P_n be n polytopes in \mathbb{R}^n and $\lambda_1, \dots, \lambda_n$ non-negative real parameters. Then $\text{vol}_n(\lambda_1 P_1 + \dots + \lambda_n P_n)$ is a homogeneous polynomial of degree n in $\lambda_1, \dots, \lambda_n$ with non-negative coefficients. The coefficient of $\lambda_1 \cdots \lambda_n$ is called the **mixed volume** of P_1, \dots, P_n and is denoted by $\text{MV}_n(P_1, \dots, P_n)$.

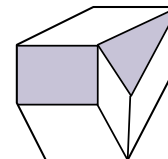
$$\text{MV}_n(P_1, \dots, P_n) = \sum_{\substack{C \text{ mixed cell of a} \\ \text{mixed subdivision} \\ \text{of } P_1 + \dots + P_n}} \text{vol}_n(C)$$

- Symmetry: $\text{MV}_n(P_1, \dots, P_n) = \text{MV}_n(P_{\sigma(1)}, \dots, P_{\sigma(n)})$ for a permutation σ .
- Linearity: $\text{MV}_n(\alpha P_1 + \beta P'_1, \dots, P_n) = \alpha \text{MV}_n(P_1, \dots, P_n) + \beta \text{MV}_n(P'_1, \dots, P_n)$
- Generalizes the Volume: $\text{MV}_n(P, \dots, P) = n! \text{vol}_n(P)$
- Invariance under rigid motions: $\text{MV}_n(P_1, \dots, P_n) = \text{MV}_n(\alpha(P_1), \dots, \alpha(P_n))$ for any volume preserving α .



Constructing Mixed Subdivisions

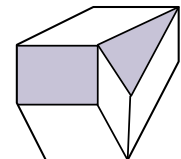
To construct a mixed subdivision of $P_1 + \dots + P_n$ we proceed as follows.



Constructing Mixed Subdivisions

To construct a mixed subdivision of $P_1 + \dots + P_n$ we proceed as follows.

- Choose a linear lifting function $\mu_i : \mathbb{R}^n \rightarrow \mathbb{R}$ for each P_i .
- Lift the polytopes P_i to $\hat{P}_i := \{(q, \mu_i(q)) : q \in P_i\} \subset \mathbb{R}^{n+1}$.
- Compute the lower hull of $\hat{P} := \hat{P}_1 + \dots + \hat{P}_n$, i.e. those facets of \hat{P} which have an inner normal with positive $(n+1)$ -coordinate.
- Project the lower hull back to \mathbb{R}^n by forgetting the last coordinate.



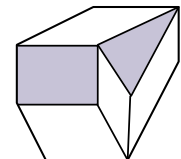
Constructing Mixed Subdivisions

To construct a mixed subdivision of $P_1 + \dots + P_n$ we proceed as follows.

- Choose a linear lifting function $\mu_i : \mathbb{R}^n \rightarrow \mathbb{R}$ for each P_i .
- Lift the polytopes P_i to $\hat{P}_i := \{(q, \mu_i(q)) : q \in P_i\} \subset \mathbb{R}^{n+1}$.
- Compute the lower hull of $\hat{P} := \hat{P}_1 + \dots + \hat{P}_n$, i.e. those facets of \hat{P} which have an inner normal with positive $(n+1)$ -coordinate.
- Project the lower hull back to \mathbb{R}^n by forgetting the last coordinate.

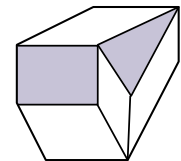
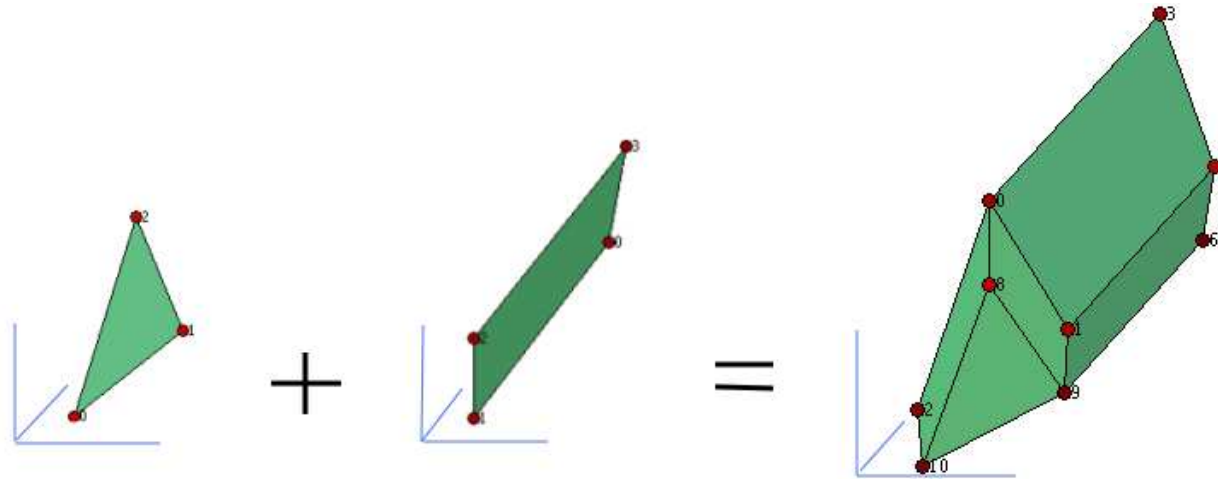
We call such a subdivision **coherent** and we will say it is **induced by μ_1, \dots, μ_n** .

A subdivision induced by μ_1, \dots, μ_n is mixed if every vertex of the lower hull can be expressed uniquely as a Minkowski sum. Such a set of liftings will be called **(sufficiently) generic**.



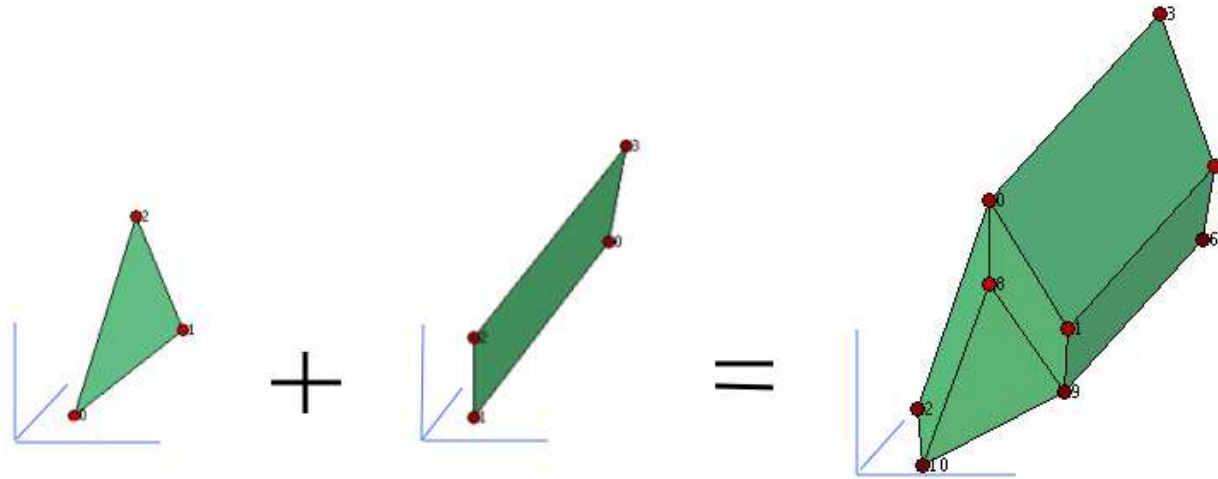
Constructing Mixed Subdivisions – Example

Lifting the rectangle and the triangle by two different linear functions yields the following 3 dimensional objects.

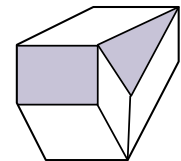
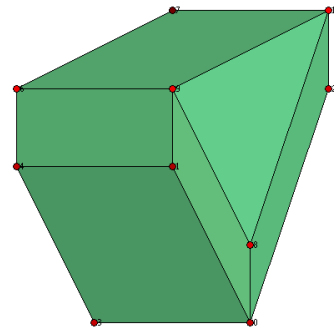


Constructing Mixed Subdivisions – Example

Lifting the rectangle and the triangle by two different linear functions yields the following 3 dimensional objects.



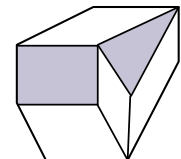
The lower hull of this Minkowski sum looks like this.



Lemma. *Let P_1, \dots, P_k be polytopes in \mathbb{R}^{m+k} and Q_1, \dots, Q_m be polytopes in $\mathbb{R}^m \subset \mathbb{R}^{m+k}$. Then*

$$\text{MV}_{m+k}(Q_1, \dots, Q_m, P_1, \dots, P_k) = \text{MV}_m(Q_1, \dots, Q_m) * \text{MV}_k(\pi(P_1), \dots, \pi(P_k))$$

where $\pi : \mathbb{R}^{m+k} \rightarrow \mathbb{R}^k$ denotes the projection on the last k coordinates.



Lemma. Let P_1, \dots, P_k be polytopes in \mathbb{R}^{m+k} and Q_1, \dots, Q_m be polytopes in $\mathbb{R}^m \subset \mathbb{R}^{m+k}$. Then

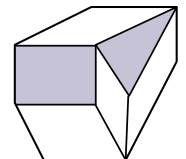
$$\text{MV}_{m+k}(Q_1, \dots, Q_m, P_1, \dots, P_k) = \text{MV}_m(Q_1, \dots, Q_m) * \text{MV}_k(\pi(P_1), \dots, \pi(P_k))$$

where $\pi : \mathbb{R}^{m+k} \rightarrow \mathbb{R}^k$ denotes the projection on the last k coordinates.

Proof: Assume that $P_1 = \dots = P_k = P$ and $Q_1 = \dots = Q_m = Q$. Then we have to show that

$$\text{MV}_{m+k}(Q, \dots, Q, P, \dots, P) = m! k! * \text{vol}_m(Q) * \text{vol}_k(\pi(P)) .$$

But this is a known result (Betke, Ewald).



Lemma. *Let P_1, \dots, P_k be polytopes in \mathbb{R}^{m+k} and Q_1, \dots, Q_m be polytopes in $\mathbb{R}^m \subset \mathbb{R}^{m+k}$. Then*

$$\text{MV}_{m+k}(Q_1, \dots, Q_m, P_1, \dots, P_k) = \text{MV}_m(Q_1, \dots, Q_m) * \text{MV}_k(\pi(P_1), \dots, \pi(P_k))$$

where $\pi : \mathbb{R}^{m+k} \rightarrow \mathbb{R}^k$ denotes the projection on the last k coordinates.

Proof: Assume that $P_1 = \dots = P_k = P$ and $Q_1 = \dots = Q_m = Q$. Then we have to show that

$$\text{MV}_{m+k}(Q, \dots, Q, P, \dots, P) = m! k! * \text{vol}_m(Q) * \text{vol}_k(\pi(P)) .$$

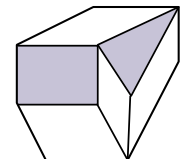
But this is a known result (Betke, Ewald).

We define g_1 and g_2 via

$$g_1(Q_1, \dots, Q_m, P_1, \dots, P_k) := \text{MV}_{m+k}(Q_1, \dots, Q_m, P_1, \dots, P_k)$$

$$g_2(Q_1, \dots, Q_m, P_1, \dots, P_k) := \text{MV}_m(Q_1, \dots, Q_m) * \text{MV}_k(\pi(P_1), \dots, \pi(P_k)) .$$

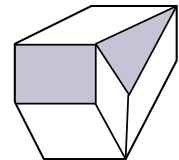
Both are invariant under reordering the Q_i or the P_j and linear in each argument.



The Separation Lemma: Sketch of the Proof 1

Let $f : A \times \dots \times A \rightarrow B$ be a symmetric multilinear function, where A and B are semigroups. By expanding the right hand side it can be seen that

$$f(a_1, \dots, a_n) = \frac{1}{n!} \sum_{1 \leq i_1 < \dots < i_q \leq n} (-1)^{n-q} f(a_{i_1} + \dots + a_{i_q}, \dots, a_{i_1} + \dots + a_{i_q}) .$$



The Separation Lemma: Sketch of the Proof 1

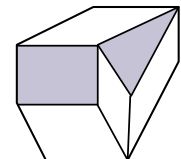
Let $f : A \times \dots \times A \rightarrow B$ be a symmetric multilinear function, where A and B are semigroups. By expanding the right hand side it can be seen that

$$f(a_1, \dots, a_n) = \frac{1}{n!} \sum_{1 \leq i_1 < \dots < i_q \leq n} (-1)^{n-q} f(a_{i_1} + \dots + a_{i_q}, \dots, a_{i_1} + \dots + a_{i_q}) .$$

The functions

$$\begin{aligned} \tilde{g}_i^{(P_1, \dots, P_k)}(Q_1, \dots, Q_m) &:= g_i(Q_1, \dots, Q_m, P_1, \dots, P_k) \text{ and} \\ \bar{g}_i^{(Q)}(P_1, \dots, P_k) &:= g_i(Q, \dots, Q, P_1, \dots, P_k) \text{ for } i = 1, 2 \end{aligned}$$

satisfy these conditions.



The Separation Lemma: Sketch of the Proof 1

Let $f : A \times \dots \times A \rightarrow B$ be a symmetric multilinear function, where A and B are semigroups. By expanding the right hand side it can be seen that

$$f(a_1, \dots, a_n) = \frac{1}{n!} \sum_{1 \leq i_1 < \dots < i_q \leq n} (-1)^{n-q} f(a_{i_1} + \dots + a_{i_q}, \dots, a_{i_1} + \dots + a_{i_q}) .$$

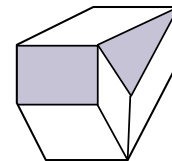
The functions

$$\begin{aligned} \tilde{g}_i^{(P_1, \dots, P_k)}(Q_1, \dots, Q_m) &:= g_i(Q_1, \dots, Q_m, P_1, \dots, P_k) \text{ and} \\ \bar{g}_i^{(Q)}(P_1, \dots, P_k) &:= g_i(Q, \dots, Q, P_1, \dots, P_k) \text{ for } i = 1, 2 \end{aligned}$$

satisfy these conditions.

Hence for $i = 1, 2$ we can expand $g_i(Q_1, \dots, Q_m, P_1, \dots, P_k)$ twice and we see that both functions g_1 and g_2 are fully determined by their images on tuples of polytopes where $Q_1 = \dots = Q_m = Q$ and $P_1 = \dots = P_k = P$. This proves the Lemma.

Remark: Equiv. result probably know: V.P. Fedotov. The sum of p-th surface functions. (Ukrain. Geom. Sb.)



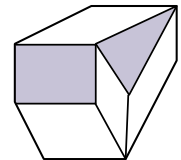
Lemma. *Given polytopes $P_1, \dots, P_n \subset \mathbb{R}^n$ and lifting vectors $\mu_1, \dots, \mu_n \in \mathbb{R}_{\geq 0}^n$. Denote the vertices of P_i by $v_1^{(i)}, \dots, v_{r_i}^{(i)}$ and choose one edge $e_i = \{v_{k_i}^{(i)}, v_{l_i}^{(i)}\}$ from each P_i . Then $C := e_1 + \dots + e_n$ is a mixed cell of the mixed subdivision induced by the liftings μ_i if and only if*

i) The edge matrix $E := V_a - V_b$ is non-singular (where $V_a := (v_{k_1}^{(1)}, \dots, v_{k_n}^{(n)})$ and $V_b := (v_{l_1}^{(1)}, \dots, v_{l_n}^{(n)})$) and

ii) For all polytopes P_i and all vertices $v_s^{(i)}$ of P_i which are not in e_i we have:

$$\left(\text{diag}(\mu^T E)^T E^{-1} - \mu_i^T \right) \cdot \left(v_{l_i}^{(i)} - v_s^{(i)} \right) \geq 0$$

where $\mu := (\mu_1, \dots, \mu_n)$ and where $\text{diag}(V)$ denotes the vector of the diagonal entries of V .



Lemma. Given polytopes $P_1, \dots, P_n \subset \mathbb{R}^n$ and lifting vectors $\mu_1, \dots, \mu_n \in \mathbb{R}_{\geq 0}^n$. Denote the vertices of P_i by $v_1^{(i)}, \dots, v_{r_i}^{(i)}$ and choose one edge $e_i = \{v_{k_i}^{(i)}, v_{l_i}^{(i)}\}$ from each P_i . Then $C := e_1 + \dots + e_n$ is a mixed cell of the mixed subdivision induced by the liftings μ_i if and only if

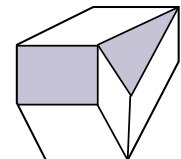
i) The edge matrix $E := V_a - V_b$ is non-singular (where $V_a := (v_{k_1}^{(1)}, \dots, v_{k_n}^{(n)})$ and $V_b := (v_{l_1}^{(1)}, \dots, v_{l_n}^{(n)})$) and

ii) For all polytopes P_i and all vertices $v_s^{(i)}$ of P_i which are not in e_i we have:

$$\left(\text{diag}(\mu^T E)^T E^{-1} - \mu_i^T \right) \cdot \left(v_{l_i}^{(i)} - v_s^{(i)} \right) \geq 0$$

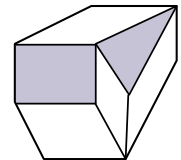
where $\mu := (\mu_1, \dots, \mu_n)$ and where $\text{diag}(V)$ denotes the vector of the diagonal entries of V .

Proof (Sketch). Write the condition that C is a mixed cell as a linear program and then use linear programming duality to get explicit conditions for the optimality and hence for C to be a mixed cell.



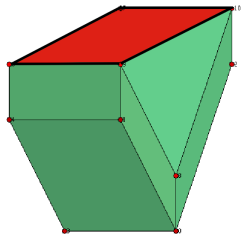
The Complete Fan of Mixed Subdivisions I

We demonstrate this with a 2-dimensional example. Which liftings μ_1 and μ_2 induce a subdivision that has the red cell as a mixed cell?



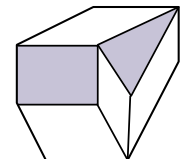
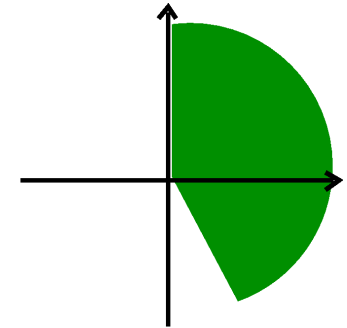
The Complete Fan of Mixed Subdivisions I

We demonstrate this with a 2-dimensional example. Which liftings μ_1 and μ_2 induce a subdivision that has the red cell as a mixed cell?



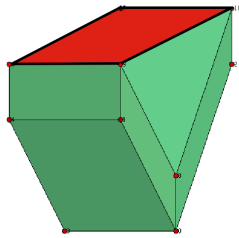
$$C = \left\{ \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix} \right\} + \left\{ \begin{pmatrix} 0 \\ 2 \end{pmatrix}, \begin{pmatrix} 2 \\ 3 \end{pmatrix} \right\}$$

Lemma $\xrightarrow{\quad} \begin{pmatrix} 2 \\ 1 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$ and $\begin{pmatrix} 1 \\ 0 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$



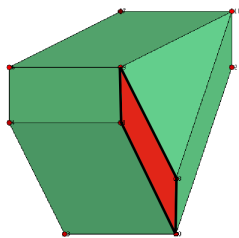
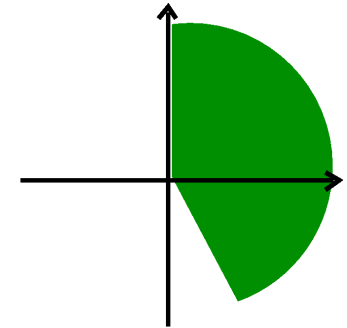
The Complete Fan of Mixed Subdivisions I

We demonstrate this with a 2-dimensional example. Which liftings μ_1 and μ_2 induce a subdivision that has the red cell as a mixed cell?



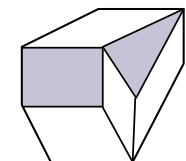
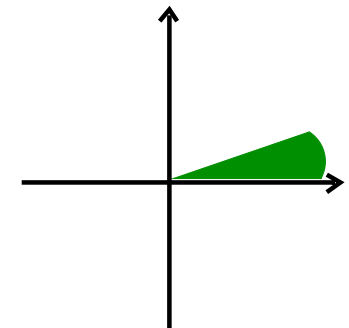
$$C = \left\{ \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix} \right\} + \left\{ \begin{pmatrix} 0 \\ 2 \end{pmatrix}, \begin{pmatrix} 2 \\ 3 \end{pmatrix} \right\}$$

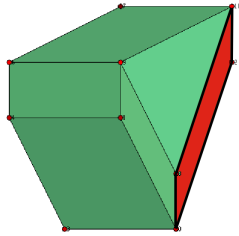
Lemma $\rightarrow \begin{pmatrix} 2 \\ 1 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$ and $\begin{pmatrix} 1 \\ 0 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$



$$C = \left\{ \begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix} \right\} + \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 2 \end{pmatrix} \right\}$$

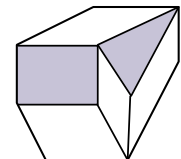
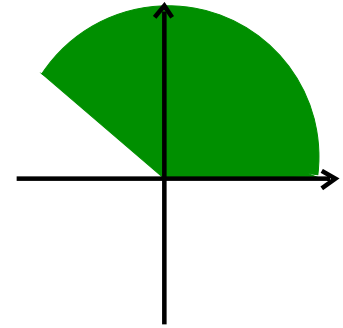
Lemma $\rightarrow \begin{pmatrix} 1 \\ -2 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$ and $\begin{pmatrix} 0 \\ 1 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$



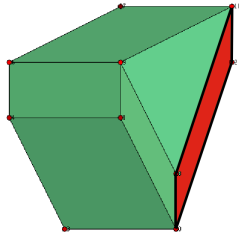


$$C = \left\{ \begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix} \right\} + \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 3 \end{pmatrix} \right\}$$

Lemma $\longrightarrow \begin{pmatrix} 1 \\ 3 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$ and $\begin{pmatrix} 0 \\ 1 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$

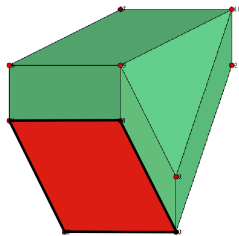
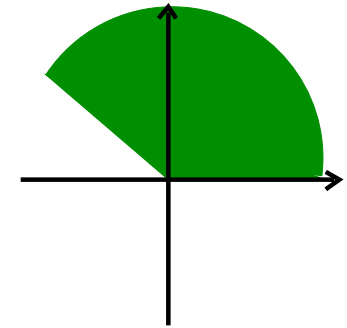


The Complete Fan of Mixed Subdivisions II



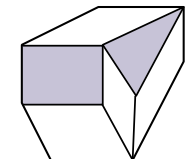
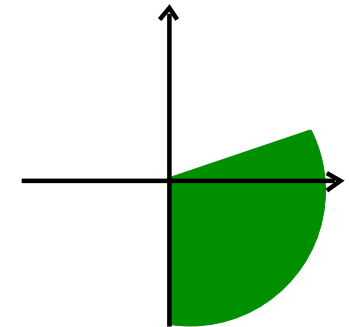
$$C = \left\{ \begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix} \right\} + \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 3 \end{pmatrix} \right\}$$

Lemma $\longrightarrow \begin{pmatrix} 1 \\ 3 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$ and $\begin{pmatrix} 0 \\ 1 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$

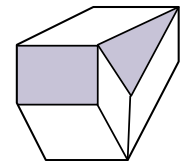
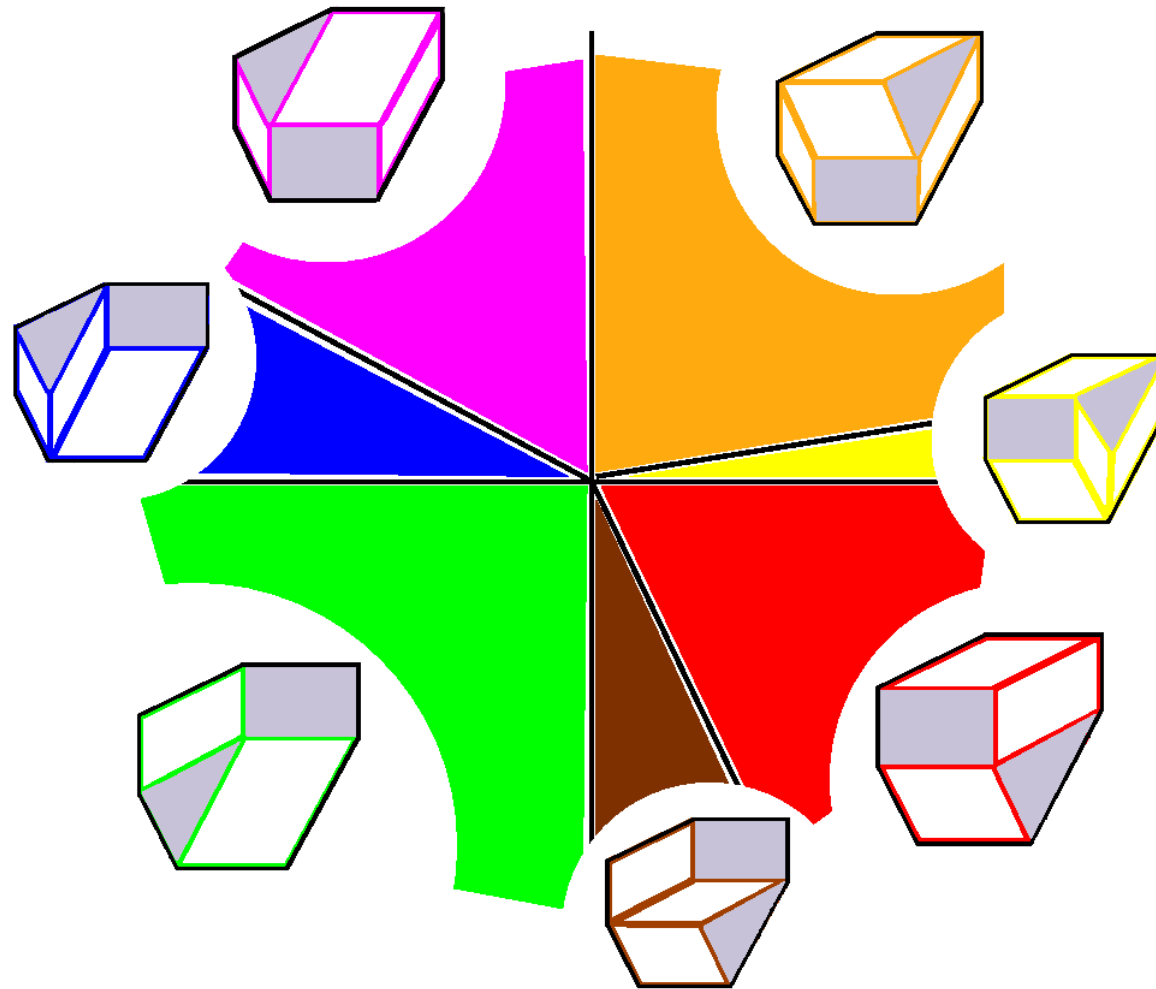


$$C = \left\{ \begin{pmatrix} 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 2 \\ 0 \end{pmatrix} \right\} + \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 2 \end{pmatrix} \right\}$$

Lemma $\longrightarrow \begin{pmatrix} 1 \\ -2 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$ and $\begin{pmatrix} 1 \\ 0 \end{pmatrix} (\mu_2 - \mu_1) \geq 0$



The Complete Fan of Mixed Subdivisions – Picture

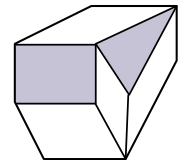


Bernstein's Theorem

Theorem (Bernstein). *Given polynomials f_1, \dots, f_n over \mathbb{C} with finitely many common zeroes in $(\mathbb{C}^*)^n$, let P_i denote the Newton polytope of f_i in \mathbb{R}^n .*

Then the number of common zeroes of the f_i in $(\mathbb{C}^)^n$ is bounded above by the mixed volume $MV_n(P_1, \dots, P_n)$.*

Moreover for generic choices of the coefficients in the f_i , the number of common solutions is exactly $MV_n(P_1, \dots, P_n)$.



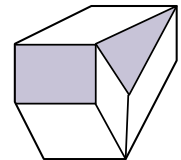
Bernstein's Theorem

Theorem (Bernstein). *Given polynomials f_1, \dots, f_n over \mathbb{C} with finitely many common zeroes in $(\mathbb{C}^*)^n$, let P_i denote the Newton polytope of f_i in \mathbb{R}^n .*

Then the number of common zeroes of the f_i in $(\mathbb{C}^)^n$ is bounded above by the mixed volume $MV_n(P_1, \dots, P_n)$.*

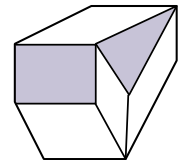
Moreover for generic choices of the coefficients in the f_i , the number of common solutions is exactly $MV_n(P_1, \dots, P_n)$.

- Bernstein also gives explicit conditions when coefficients are generic or not.
- The bound on the number of solutions obtained from the mixed volume is called the **BKK-bound** due to the work of Bernstein, Kushnirenko and Khovanskii.
- An upper bound for the common solutions in \mathbb{C} for the system above is $MV_n(\text{conv}(P_1 \cup 0), \dots, \text{conv}(P_n \cup 0))$ (Li, Wang).
- The BKK-bound generalizes the Bézout bound.



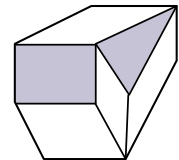
Laman Graphs

- A **Laman Graph** is a graph $G = (V, E)$ with $|E| = 2|V| - 3$ edges where each subset of k vertices spans at most $2k - 3$ edges.
- A **framework** is a tuple (G, L) where $G = (V, E)$ is a graph and $L = \{l_{i,j} : [v_i, v_j] \in E\}$ is a set of $|E|$ positive numbers interpreted as edge lengths.



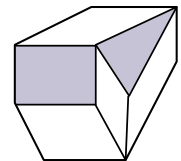
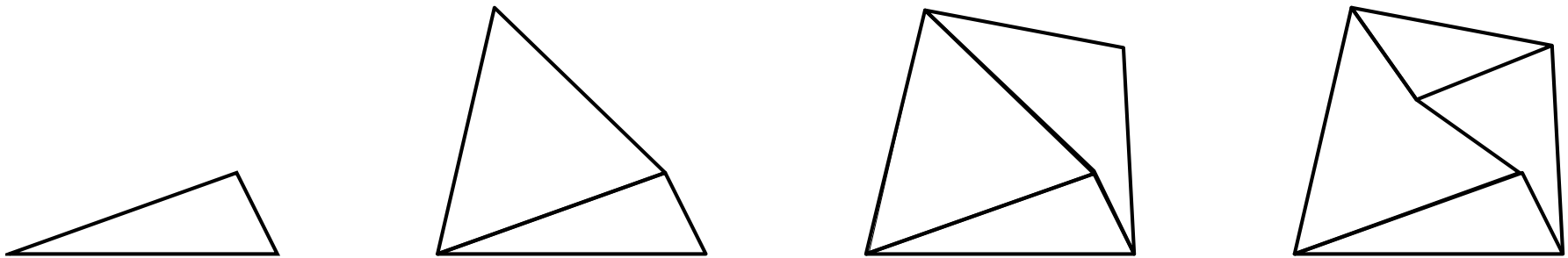
Laman Graphs

- A **Laman Graph** is a graph $G = (V, E)$ with $|E| = 2|V| - 3$ edges where each subset of k vertices spans at most $2k - 3$ edges.
- A **framework** is a tuple (G, L) where $G = (V, E)$ is a graph and $L = \{l_{i,j} : [v_i, v_j] \in E\}$ is a set of $|E|$ positive numbers interpreted as edge lengths.
- For generic edge lengths, Laman graphs are minimally rigid.



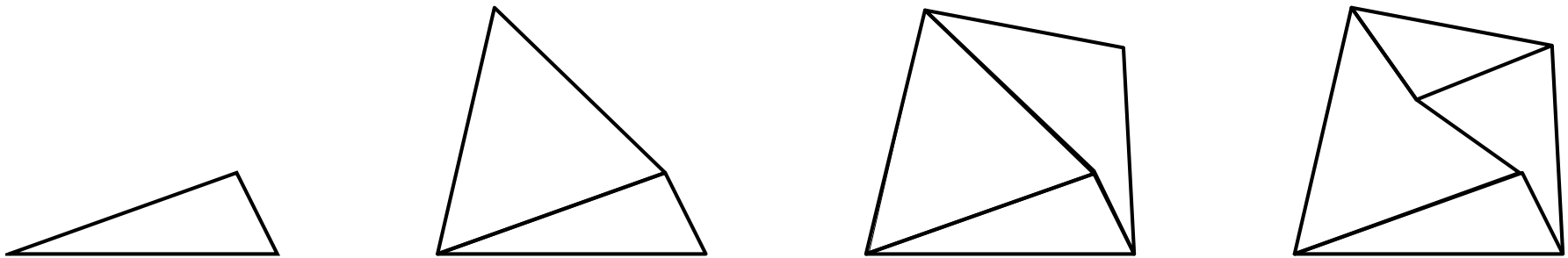
Laman Graphs

- A **Laman Graph** is a graph $G = (V, E)$ with $|E| = 2|V| - 3$ edges where each subset of k vertices spans at most $2k - 3$ edges.
- A **framework** is a tuple (G, L) where $G = (V, E)$ is a graph and $L = \{l_{i,j} : [v_i, v_j] \in E\}$ is a set of $|E|$ positive numbers interpreted as edge lengths.
- For generic edge lengths, Laman graphs are minimally rigid.
- Each Laman Graph can be constructed via a Henneberg sequence.

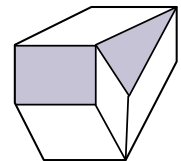


Laman Graphs

- A **Laman Graph** is a graph $G = (V, E)$ with $|E| = 2|V| - 3$ edges where each subset of k vertices spans at most $2k - 3$ edges.
- A **framework** is a tuple (G, L) where $G = (V, E)$ is a graph and $L = \{l_{i,j} : [v_i, v_j] \in E\}$ is a set of $|E|$ positive numbers interpreted as edge lengths.
- For generic edge lengths, Laman graphs are minimally rigid.
- Each Laman Graph can be constructed via a Henneberg sequence.

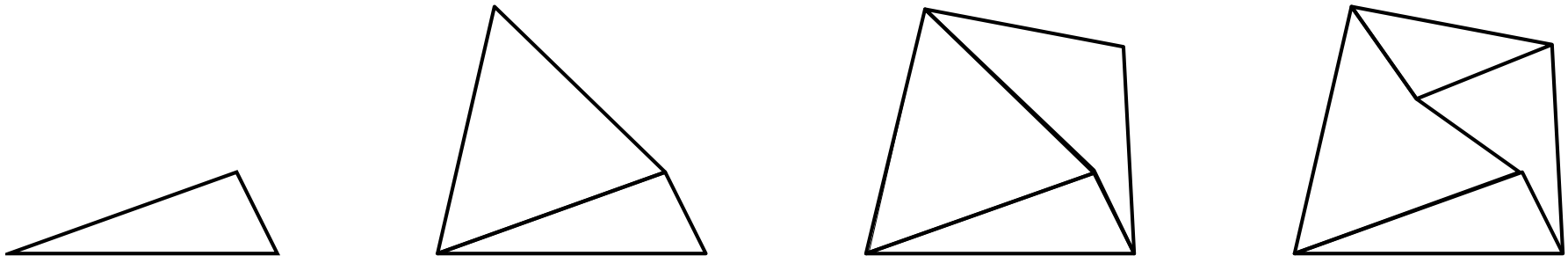


A **Henneberg I step** adds one new vertex v and two new edges, connecting v to two arbitrary previous vertices .



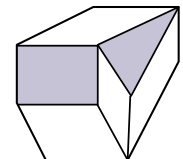
Laman Graphs

- A **Laman Graph** is a graph $G = (V, E)$ with $|E| = 2|V| - 3$ edges where each subset of k vertices spans at most $2k - 3$ edges.
- A **framework** is a tuple (G, L) where $G = (V, E)$ is a graph and $L = \{l_{i,j} : [v_i, v_j] \in E\}$ is a set of $|E|$ positive numbers interpreted as edge lengths.
- For generic edge lengths, Laman graphs are minimally rigid.
- Each Laman Graph can be constructed via a Henneberg sequence.



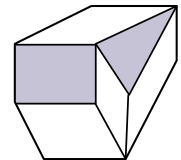
A **Henneberg I step** adds one new vertex v and two new edges, connecting v to two arbitrary previous vertices .

A **Henneberg II step** adds one new vertex v and three new edges, connecting v to three previous vertices such that at least two of these vertices are connected via an edge e and this certain edge is removed.



The Problem and its Formulation

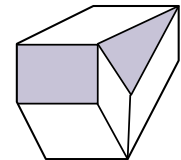
Question: Given a Laman graph on $|V|$ vertices with generic edge lengths, how many embeddings does it have modulo rigid motions (translations and rotation)?



The Problem and its Formulation

Question: Given a Laman graph on $|V|$ vertices with generic edge lengths, how many embeddings does it have modulo rigid motions (translations and rotation)?

- The best bound known so far is $\binom{2|V|-4}{|V|-2}$ (Borcea and Streinu '04).
- We have an upper bound of order $4^{|V|}$ and a lower bound of order $2.28^{|V|}$.

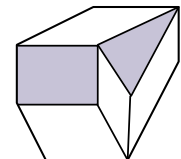


The Problem and its Formulation

Question: Given a Laman graph on $|V|$ vertices with generic edge lengths, how many embeddings does it have modulo rigid motions (translations and rotation)?

- The best bound known so far is $\binom{2|V|-4}{|V|-2}$ (Borcea and Streinu '04).
- We have an upper bound of order $4^{|V|}$ and a lower bound of order $2.28^{|V|}$.

Each prescribed edge length translates into a polynomial equation. I.e. if $[v_i, v_j] \in E$ with length $L_{i,j}$, we require $(x_i - x_j)^2 + (y_i - y_j)^2 = L_{i,j}^2$ where (x_i, y_i) is the embedding of v_i .



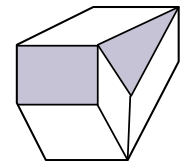
The Problem and its Formulation

Question: Given a Laman graph on $|V|$ vertices with generic edge lengths, how many embeddings does it have modulo rigid motions (translations and rotation)?

- The best bound known so far is $\binom{2|V|-4}{|V|-2}$ (Borcea and Streinu '04).
- We have an upper bound of order $4^{|V|}$ and a lower bound of order $2.28^{|V|}$.

Each prescribed edge length translates into a polynomial equation. I.e. if $[v_i, v_j] \in E$ with length $L_{i,j}$, we require $(x_i - x_j)^2 + (y_i - y_j)^2 = L_{i,j}^2$ where (x_i, y_i) is the embedding of v_i .

To get rid of translations and rotations we fix one point $(x_1, y_1) = (c_1, c_2)$ and the direction of the edge $[v_1, v_2]$ by setting $y_2 = c_3$.



Question: Given a Laman graph on $|V|$ vertices with generic edge lengths, how many embeddings does it have modulo rigid motions (translations and rotation)?

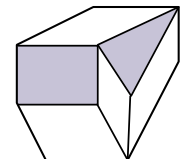
- The best bound known so far is $\binom{2|V|-4}{|V|-2}$ (Borcea and Streinu '04).
- We have an upper bound of order $4^{|V|}$ and a lower bound of order $2.28^{|V|}$.

Each prescribed edge length translates into a polynomial equation. I.e. if $[v_i, v_j] \in E$ with length $L_{i,j}$, we require $(x_i - x_j)^2 + (y_i - y_j)^2 = L_{i,j}^2$ where (x_i, y_i) is the embedding of v_i .

To get rid of translations and rotations we fix one point $(x_1, y_1) = (c_1, c_2)$ and the direction of the edge $[v_1, v_2]$ by setting $y_2 = c_3$.

Hence we want to study the solutions to the following system.

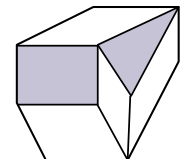
$$\left\{ \begin{array}{l} x_1 - c_1 = 0, \quad y_1 - c_2 = 0, \quad x_2 - (L_{1,2} - c_1) = 0, \quad y_2 - c_3 = 0 \\ (x_i - x_j)^2 + (y_i - y_j)^2 - L_{i,j}^2 = 0 \quad \forall [v_i, v_j] \in E - \{[v_1, v_2]\} \end{array} \right\}$$



Theorem. *The mixed volume of our initial system*

$$\left\{ \begin{array}{l} x_1 - c_1 = 0, \quad y_1 - c_2 = 0 \\ x_2 - (L_{1,2} - c_1) = 0, \quad y_2 - c_3 = 0 \\ (x_i - x_j)^2 + (y_i - y_j)^2 - L_{i,j}^2 = 0 \\ \quad \forall [v_i, v_j] \in E - \{[v_1, v_2]\} \end{array} \right\}$$

is exactly $4^{|V|-2}$.

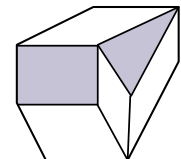


Theorem. *The mixed volume of our initial system*

$$\left\{ \begin{array}{l} x_1 - c_1 = 0, \quad y_1 - c_2 = 0 \\ x_2 - (L_{1,2} - c_1) = 0, \quad y_2 - c_3 = 0 \\ (x_i - x_j)^2 + (y_i - y_j)^2 - L_{i,j}^2 = 0 \\ \forall [v_i, v_j] \in E - \{[v_1, v_2]\} \end{array} \right\}$$

is exactly $4^{|V|-2}$.

- Proof (Sketch).**
- Bound the mixed volume from above with the product of the degrees.
 - Using the second Lemma we find lifting vectors that induce "large" mixed cells such that the volume adds up to this value.



Theorem. *The mixed volume of our initial system*

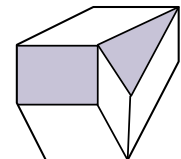
$$\left\{ \begin{array}{l} x_1 - c_1 = 0, \quad y_1 - c_2 = 0 \\ x_2 - (L_{1,2} - c_1) = 0, \quad y_2 - c_3 = 0 \\ (x_i - x_j)^2 + (y_i - y_j)^2 - L_{i,j}^2 = 0 \\ \forall [v_i, v_j] \in E - \{[v_1, v_2]\} \end{array} \right\}$$

is exactly $4^{|V|-2}$.

Proof (Sketch). • Bound the mixed volume from above with the product of the degrees.

- Using the second Lemma we find lifting vectors that induce "large" mixed cells such that the volume adds up to this value.

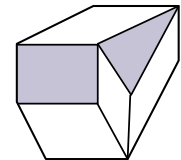
Corollary. *The number of embeddings of a Laman graph framework with generic edge lengths is strictly less than $4^{|V|-2}$.*



Theorem. *A Henneberg I step **at most doubles** the number of embeddings of the framework and there is always a choice of edge lengths such that the number of embeddings is doubled.*

Corollary. *The number of embeddings of Henneberg I graphs is less than or equal $2^{|V|-2}$ and this bound is sharp.*

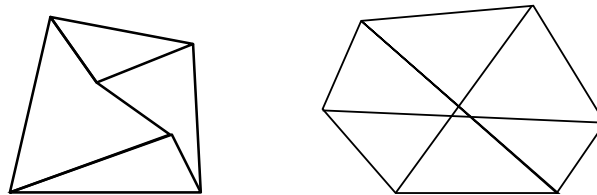
The only Henneberg II Laman graphs on 6 vertices are the Desargues graph and $K_{3,3}$.



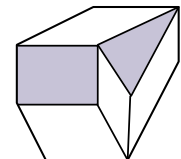
Theorem. *A Henneberg I step **at most doubles** the number of embeddings of the framework and there is always a choice of edge lengths such that the number of embeddings is doubled.*

Corollary. *The number of embeddings of Henneberg I graphs is less than or equal $2^{|V|-2}$ and this bound is sharp.*

The only Henneberg II Laman graphs on 6 vertices are the Desargues graph and $K_{3,3}$.



- The Mixed Volume bound is **32**. (Compared to 70 from the Borcea-Streinu bound.)
- For the Desargues graph Borcea and Streinu show that the number of embeddings is exactly 24.
- For $K_{3,3}$ Husty showed that there are at least 32 embeddings.



The End

Danke!

